

# IMS Presence Server: Traffic Analysis & Performance Modelling



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**Presented** 

By

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- Background Introduction
- Presence Server Traffic Load Analysis
- Traffic Process Modelling of a Presence Server
  - An Accurate Model
  - A Simplified Model
- Conclusion

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#### **Presence Service Introduction**

- Presence is a service that allows a user to be informed about the reachability, availability, and willingness of communication of another user.
- Able to indicate users' status: online or not; Idle or busy; communication means and capabilities: SMS, MMS, 3G,2G Phone....
- A key enabler for many popular applications: push-to-talk (PTT), instant messaging (IM).
- 3G IP Multimedia Subsystem (IMS) already has presence service well supported in its architecture.



#### SIP Presence Architecture

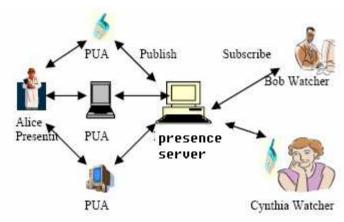
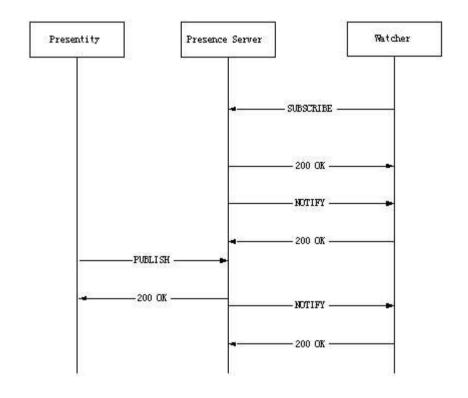


Figure 1-SIP Presence Architecture

- •Alice is the publisher whose presence information is published to a presence server by PUA.
- •PUA is the presence user agent of a user.
- •Presence Agent is a presence server that is responsible for managing presence information.
- •Bob/Cynthia are watchers who subscribes to information from PS on Alice's presentity.

# SIP Presence Service Message Flow



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#### User Behavior & Traffic Characterization

#### Login and Logout

- A user's login generates an initial PUBLISH message to the PS
- Refresh PUBLISH messages are generated periodically
- Terminating PUBLISH message is sent to the PS upon the user's logout

#### Presence Subscription

- Subscription of a user's presentity results in an initial SUBSCRIBE message being sent to the PS,
- Refresh SUBSCRIBE messages will be sent to the PS periodically
- Terminating SUBSCRIBE message un-subscribes the other user's presentity.

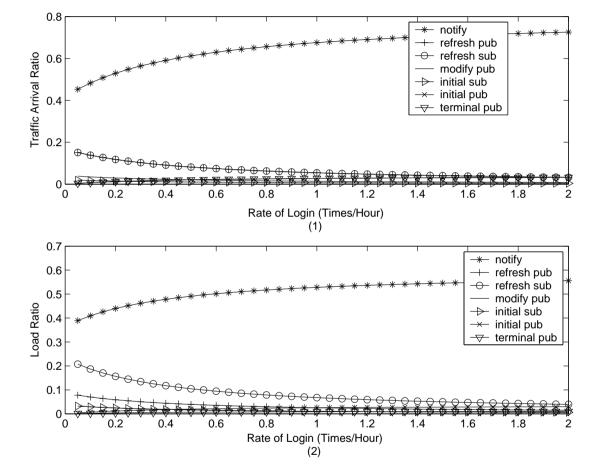
#### Presence Status Updates

• Change of a user's status results in a modifying PUBLISH message,

# Traffic Types

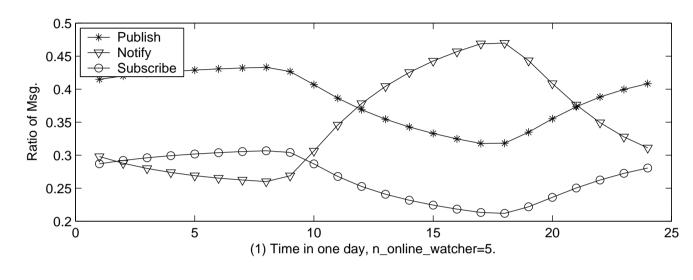
- Traffic to a PS is divided into eight types:
  - initial/refresh/modifying/terminal publish
  - initial/refresh/terminal subscribe
  - notify
- Process Time for each Traffic type is Different (a reference implementation)

#### Ratio of Traffic Load vs. User Login Rate , *online watcher* = 10

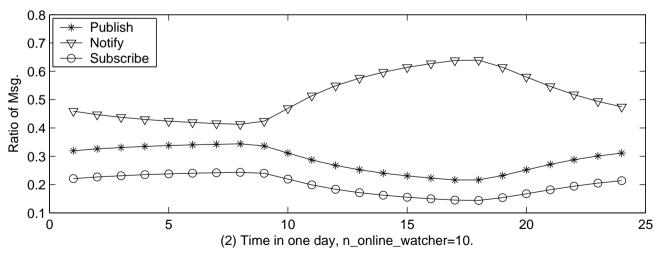


- Traffic Load increases with login frequence increases.
- Notify messages is the largest part of the traffic load.
- Refresh PUBLISH
   and refresh
   Subscribe has the
   same traffic arrival
   rate, but Refresh
   Subscribe has more
   heavier load.

# Messages Ratio Distribution During the Time of a Day.



- •Traffic to a PS varies during the time of a day
- Notify messages are the largest part of all the traffic load when online watcher number greater than 10



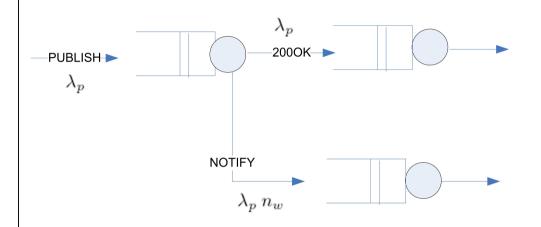
#### Presence Server Traffic Characterization

- NOTIFY messages have great impact on traffic load to a PS
- Traffic rate to a PS varies greatly with the time of a day.
- With each message to/from PS will be relayed by IMS Core network, NOTIFY message process will be critical for whole IMS network.
- NOTIFY message process mechanism in a PS is studied in the following.



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# Publish and Notify Queues in a PS

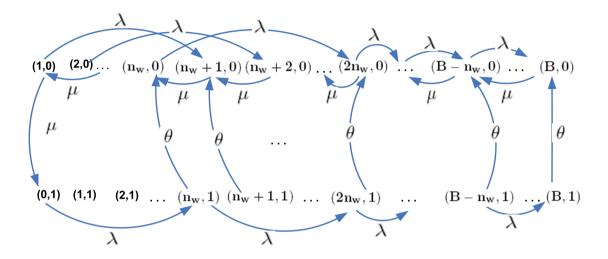


- •A PUBLISH message arrives a PS,  $n_w$  NOTIFY messages are generated.
- •2000K messages are sent out as soon as possible to prevent retransmission
- •NOTIFY messages are buffered in a queue and controlled by a scheduling mechanism.

Note: nw is the number of online watchers

•Notify messages queue is modelled as a Queueing System with Controlled Vacation and Batch Poisson Arrival

# State Transition Graph of NOTIFY Queue



•Based on the state transition graph, balance equations can be used to get probability of system states and the relationship between vacation time  $(T = \frac{1}{\theta})$  and state probability.

# Optimization Problem from the Analysis

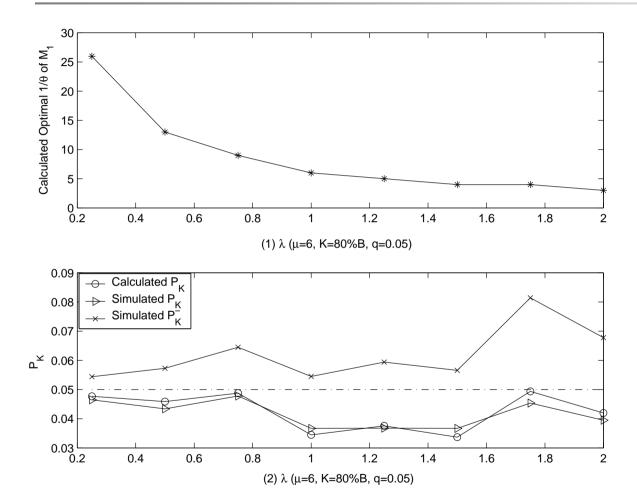
$$\min \theta$$
 (34)

s.t.

$$P_K \le q$$
 (35)

- Minimize the busy time of the Presence Server such that
- ullet The probability of message queue length greater than K should be less than q,

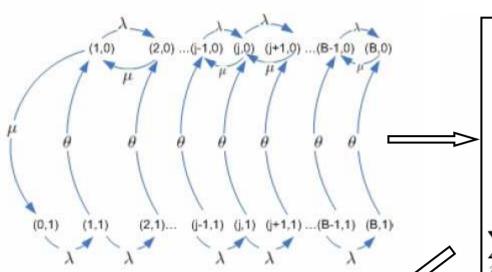
# **Optimal Notify Timers**



- •A higher arrival rate needs a shorter vacation time.
- •Timer is not very sensitive to the traffic arrival rate
- •Traffic load can be divided into several levels and each level is associated with a control timer value.
- •Optimal timer value, the loss probability of NOTIFY queue can vary for difference traffic load as indicated in Fig. 10(2), but they all satisfies the constraint.

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# A Queue System with Controlled Vacation and Poisson Arrival (A Simplified Model)



$$\lambda p_{0,1} = \mu p_{1,0}$$

$$(\lambda + \mu) p_{1,0} = \mu p_{2,0} + \theta p_{1,1}$$

$$(\lambda + \mu) p_{j,0} = \lambda p_{j-1,0} + \mu p_{j+1,0} + \theta p_{j,1},$$

$$j = 2, \dots, B - 1.$$

$$\mu p_{B,0} = \lambda p_{B-1,0} + \theta p_{B,1}$$

$$(\theta + \lambda) p_{j,1} = \lambda p_{j-1,1}, j = 1, \dots, B - 1.$$

$$\theta p_{B,1} = \lambda p_{B-1,1}$$

$$\sum_{j=1}^{B} p_{j,0} + \sum_{j=0}^{B} p_{j,1} = 1$$

$$p_{B,1} = \frac{\mu - \lambda}{\mu (1 + \frac{\theta}{\lambda})^B + \mu \times J - \lambda \times (1 + \frac{\theta}{\lambda})^B - I\lambda}$$

$$I = \frac{\theta}{\lambda} \left( 1 + \frac{\theta}{\lambda} \right)^{B-1} \frac{\lambda \left( 1 - (\frac{\lambda}{\mu})^B \right)}{\mu - \lambda}$$

$$- \frac{\theta}{\mu} \frac{\theta}{\lambda} \sum_{j=0}^{B-2} \left( \frac{\lambda}{\mu} \right)^j \sum_{k=1}^{B-j-1} \left( 1 + \frac{\theta}{\lambda} \right)^{B-k-1}$$

$$J = \frac{\lambda (B-1)}{\mu} \frac{\theta}{\lambda} \left( 1 + \frac{\theta}{\lambda} \right)^{B-1}$$

$$- \frac{\theta}{\mu} \sum_{j=1}^{B-1} (B-j) \frac{\theta}{\lambda} \left( 1 + \frac{\theta}{\lambda} \right)^{B-j-1}$$

$$p_{0,1} = \frac{\theta}{\lambda} \left( 1 + \frac{\theta}{\lambda} \right)^{B-1} p_{B,1}$$

$$p_{j,1} = \frac{\theta}{\lambda} \left( 1 + \frac{\theta}{\lambda} \right)^{B-1-j} p_{B,1}$$

$$j = 1, \dots, B - 1.$$

$$p_{1,0} = \frac{\lambda}{\mu} p_{0,1}$$

$$p_{j,0} = \frac{\lambda}{\mu} p_{j-1,0} + \frac{\lambda}{\mu} p_{0,1} - \frac{\theta}{\mu} \sum_{k=1}^{j-1} p_{k,1},$$

$$j = 2, \dots, B.$$



 $\min \theta$ 

s.t.

$$P_K \leq q$$

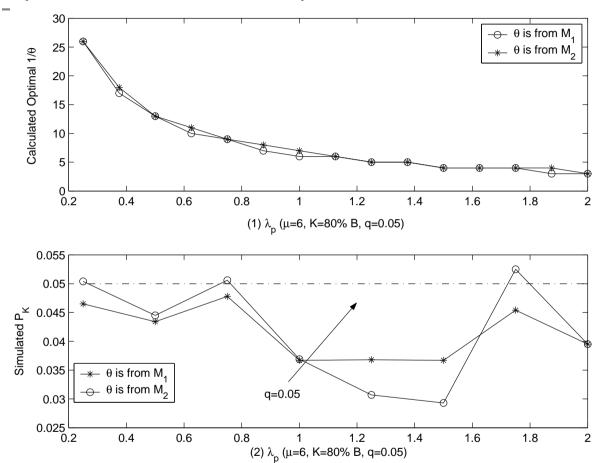
$$P_K = \sum_{j=K+1}^{B} (p_{j,1} + p_{j,0})$$

#### **Buffer Control Optimization Algorithm 1:**

#### begin

- 1. Given PUBLISH traffic arrival rate  $\lambda_p$ , number of online watchers  $n_w$ , QoS parameter K, q and timer interval I = [L, U], i = 1, maxI = L;
- 2. for i=L:U /\* loop from L to U \*/  $\theta = 1/i$ , Solve equations array  $M_1$  to get  $p_{j,0}, p_{j,1}, j = 0, \cdots, B$  Get  $P_K$  from equation  $P_K = \sum_{j=K+1}^B (p_{j,1} + p_{j,0})$  if  $P_K \leq q$  and i > maxI, maxI = i; end
- 3.  $\theta = \frac{1}{maxI}$ , Return **end**

# Experimental Result for Simplified Model



•The simplified model can satisfy the constraint at most of the time

#### **Conclusion**

- NOTIFY messages account for the largest portion of traffic load to a PS and their processing overhead has great impact on the quality of service of the PS.
- Queuing system with Controlled Vacation and Batch Poisson Arrival can be used to model Notify message queue.
- Simplified queuing system with Controlled Vacation and Poisson Arrival can be used to optimize vacation timer.
- Further work is needed to apply the analysis to the problems such as network sizing, traffic admission control.

For any questions, please contact

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# Thank you!